

On universal quadratic identities for minors of quantum matrices

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Abstract. We give a complete combinatorial characterization of homogeneous quadratic identities of “universal character” valid for minors of quantum matrices over a field. This is obtained as a consequence of a study of quantized minors of the so-called *path matrices* associated with certain planar graphs generalizing Cauchon graphs.

Keywords: quantum matrix, q -minor, Cauchon diagram, planar graph, path matrix

1 Introduction

The idea of quantization has proved its importance to bridge commutative and non-commutative versions of certain algebraic structures and promote better understanding various aspects of the latter versions. One popular structure is the *quantized coordinate ring* $\mathcal{R} = \mathcal{O}_q(\mathfrak{M}_{m,n}(\mathbb{K}))$ of $m \times n$ matrices over a field \mathbb{K} , where q is a nonzero element of \mathbb{K} , usually called the *algebra of $m \times n$ quantum matrices*. Here \mathcal{R} is generated by entries (indeterminates) of an $m \times n$ matrix X subject to Manin’s relations [11]: for $i < \ell \leq m$ and $j < k \leq n$,

$$\begin{aligned} x_{ij}x_{ik} &= qx_{ik}x_{ij}, & x_{ij}x_{\ell j} &= qx_{\ell j}x_{ij}, \\ x_{ik}x_{\ell j} &= x_{\ell j}x_{ik} & \text{and} & \quad x_{ij}x_{\ell k} - x_{\ell k}x_{ij} = (q - q^{-1})x_{ik}x_{\ell j}. \end{aligned} \tag{1.1}$$

We study quadratic identities for minors of quantum matrices, or *quantum minors*. For a discussion on aspects and applications of such identities, see e.g., [6, 7, 8, 9, 12] (where the list is incomplete). We present a novel, and rather transparent, combinatorial method which enables us to completely characterize and efficiently verify homogeneous quadratic identities of universal character that are valid for quantum minors. The identities of our interest can be written as

$$\sum (\text{sign}_i q^{\delta_i} [I_i | J_i]_q [I'_i | J'_i]_q : i = 1, \dots, N) = 0, \tag{1.2}$$

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where $\delta_i \in \mathbb{Z}$, $\text{sign}_i \in \{+, -\}$, and $[I|J]_q$ denotes the quantum minor whose rows and columns are indexed by $I \subseteq [m]$ and $J \subseteq [n]$, respectively. (Hereinafter, for a positive integer n' , we write $[n']$ for $\{1, 2, \dots, n'\}$.) The homogeneity means that each of the sets $I_i \cup I'_i$, $I_i \cap I'_i$, $J_i \cup J'_i$, $J_i \cap J'_i$ is invariant of i , and the term “universal” means that (1.2) should be valid independently of \mathbb{K}, q and a q -matrix (a matrix whose entries obey Manin’s relations and, possibly, additional ones). Note that any cortege $(I|J, I'|J')$ may be repeated in (1.2) many times.

Our approach has two sources. The first one is the *flow-matching* method elaborated in [5] to characterize quadratic identities for usual minors (viz. for $q = 1$). In that case the identities are viewed as

$$\sum (\text{sign}_i [I_i|J_i] [I'_i|J'_i] : i = 1, \dots, N) = 0. \quad (1.3)$$

In the method of [5], each cortege $S = (I|J, I'|J')$ determines a certain set $\mathcal{M}(S)$ of so-called *feasible matchings*. The main theorem in [5] asserts that (1.3) is valid (universally) if and only if the families \mathcal{I}^+ and \mathcal{I}^- of corteges S_i with $\text{sign}_i = +$ and $\text{sign}_i = -$, respectively, are *balanced*, in the sense that the total families of feasible matchings for corteges occurring in \mathcal{I}^+ and in \mathcal{I}^- are equal.

The second source is the path method due to Casteels [1, 2]. He associated with each Cauchon diagram C of size $m \times n$ (see [3]) a certain directed planar graph $G = G_C$ with $m + n$ distinguished vertices $r_1, \dots, r_m, c_1, \dots, c_n$, and considered the $m \times n$ *path matrix* $P_G = (p_{ij})$ of G . This matrix possesses three important properties. (i) It is a q -matrix, and therefore, $x_{ij} \mapsto p_{ij}$ gives a homomorphism of \mathcal{R} to the corresponding algebra generated by the p_{ij} . (ii) It admits an analog of Lindström’s Lemma [10]: for any $I \subseteq [m]$ and $J \subseteq [n]$ with $|I| = |J|$, the minor $[I|J]_q$ of P_G can be expressed via systems of *disjoint paths* from $\{r_i : i \in I\}$ to $\{c_j : j \in J\}$ in G . (iii) Using Cauchon’s Algorithm [3] interpreted in graph terms in [1, 2], one shows that if the diagram C is maximal (i.e., has no black cells), then Path_G becomes a *generic q -matrix* (see Corollary 3.2.5 in [2]).

In this work we consider a more general class of planar graphs, called *SE-graphs*; they possess the above properties (i),(ii) as well. Our goal is to characterize quadratic identities just for the class of path matrices of SE-graphs. Since this class contains a generic q -matrix, the identities are automatically valid in \mathcal{R} . As a result, we obtain necessary and sufficient conditions for the quantum version (in **Theorems 4.1** and **4.3**), namely: (1.2) is valid (universally) if and only if the families of corteges \mathcal{I}^+ and \mathcal{I}^- along with the function δ are *q -balanced*, which now means the existence of a bijection between the families of feasible matchings for \mathcal{I}^+ and \mathcal{I}^- that is agreeable with δ in a certain sense. Note also that our method of establishing or verifying one or another universal identity admits a rather transparent implementation.

The paper is organized as follows. Section 2 contains basic definitions and statements. Section 3 describes important ingredients and tools in our method: *double flows* (pairs of path systems related to corteges $(I|J, I'|J')$), *feasible matchings*, and transformations

of double flows by use of *exchange operations*. The crucial working tool exhibited here is [Corollary 3.5](#) which follows from a result on exchange operations proved in [4] (stated in [Theorem 3.4](#)). Based on these, Section 4 outlines a proof of the sufficiency: (1.2) is valid if the corresponding $\mathcal{I}^+, \mathcal{I}^-, \delta$ are q -balanced ([Theorem 4.1](#)). Also Section 4 contains an algorithm of recognizing the q -balancedness and one illustration and finishes with [Theorem 4.3](#) (without a proof) concerning the necessity of the q -balancedness.

For proofs in detail, additional results and numerous applications, see the full version posted in arXiv:[1604.00338\[math.QA\]](#).

2 Basic definitions and statements

Paths in graphs. Throughout, by a *graph* we mean a directed graph. A *path* in a graph $G = (V, E)$ is a sequence $P = (v_0, e_1, v_1, \dots, e_k, v_k)$ such that each e_i is an edge connecting the vertices v_{i-1}, v_i . An edge e_i is called *forward* if it is directed from v_{i-1} to v_i , denoted as $e_i = (v_{i-1}, v_i)$, and *backward* otherwise (when $e_i = (v_i, v_{i-1})$). The path P is called *directed* if it has no backward edge, and *simple* if all vertices v_i are different. When $k > 0$ and $v_0 = v_k$, P is called a *cycle*, and called a *simple cycle* if, in addition, v_1, \dots, v_k are different.

SE-graphs. A graph $G = (V, E)$ of this sort (also denoted as $(V, E; R, C)$) satisfies the following conditions:

(SE1) G is planar (with a fixed layout in the plane);

(SE2) G has edges of two types: *horizontal* edges, or *H-edges*, which are directed to the right, and *vertical* edges, or *V-edges*, which are directed downwards (so each edge points to either *south* or *east*, justifying the term “SE-graph”);

(SE3) G has two distinguished subsets of vertices: set $R = \{r_1, \dots, r_m\}$ of *sources* and set $C = \{c_1, \dots, c_n\}$ of *sinks*; moreover, r_1, \dots, r_m are disposed on a vertical line, in this order upwards, and c_1, \dots, c_n are disposed on a horizontal line, in this order from left to right; each vertex of G belongs to a directed path from R to C .

We denote by $W = W_G$ the set $V - (R \cup C)$ if *inner* vertices of G . We also say that G is an (m, n) SE-graph (where $m := |R|$ and $n := |C|$). An example is drawn in [Figure 1](#).

Remark 1. A representative special case is formed by the SE-graphs equivalent to *Cauchon graphs* introduced in [1] (which are associated with Cauchon diagrams [3]). In this case, $R = \{(0, i) : i \in [m]\}$, $C = \{(j, 0) : j \in [n]\}$, and $W \subseteq [n] \times [m]$. When $W = [n] \times [m]$, we refer to such a graph as the *extended* (m, n) -*grid* and denote by $\Gamma_{m,n}$.

Each inner vertex $v \in W$ of an SE-graph G is regarded as a *generator*. We assign the weight $w(e)$ to each edge $e = (u, v) \in E$ in a way similar to that for Cauchon graphs in [1], namely:

(W1) $w(e) := v$ if e is an H-edge with $u \in R$;

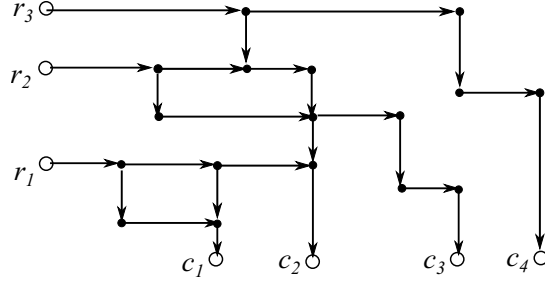


Figure 1: An SE-graph with $m = 3$ and $n = 4$

(W2) $w(e) := u^{-1}v$ if e is an H-edge and $u, v \in W$;

(W3) $w(e) := 1$ if e is a V-edge.

This gives rise to defining the weight $w(P)$ of a directed path $P = (v_0, e_1, v_1, \dots, e_k, v_k)$ to be the ordered (from left to right) product, namely:

$$w(P) := w(e_1)w(e_2) \cdots w(e_k). \quad (2.1)$$

The generators W are assumed to be subject to (quasi)commutation laws, which match those for Cauchon graphs in [1]. More precisely, for distinct $u, v \in W$,

(G1) if there is a directed *horizontal* path from u to v in G , then $uv = qvu$;

(G2) if there is a directed *vertical* path from u to v in G , then $vu = quv$;

(G3) otherwise $uv = vu$.

Quantum minors. It is convenient for us to visualize matrices in the Cartesian form: for an $m \times n$ matrix $A = (a_{ij})$, the row indices $i = 1, \dots, m$ are assumed to increase upwards, and the column indices $j = 1, \dots, n$ from left to right.

We denote by $A(I|J)$ the submatrix of A whose rows and columns are indexed by $I \subseteq [m]$ and $J \subseteq [n]$, respectively. Let $|I| = |J| =: k$, and let I consist of $i_1 < \dots < i_k$, and J consist of $j_1 < \dots < j_k$. Then the q -determinant of $A(I|J)$, or the q -minor of A for $(I|J)$, is defined as

$$[I|J]_{A,q} := \sum_{\sigma \in S_k} (-q)^{\ell(\sigma)} \prod_{d=1}^k a_{i_d j_{\sigma(d)}}, \quad (2.2)$$

where the product under \prod is ordered by increasing d , and $\ell(\sigma)$ denotes the *length* (number of inversions) of a permutation σ . The terms A and/or q in $[I|J]_{A,q}$ may be omitted when they are clear from the context.

Path matrix. An important construction in [1] associates with a Cauchon graph G a certain matrix, called the path matrix of G . This is extended to an arbitrary (m, n) SE-graph $G = (V, E)$, namely: the *path matrix* $\text{Path} = \text{Path}_G$ of G is meant to be the $m \times n$ matrix whose entries are defined by

$$\text{Path}(i|j) := \sum_{P \in \Phi_G(i|j)} w(P), \quad (i, j) \in [m] \times [n], \quad (2.3)$$

where $\Phi_G(i|j)$ is the set of directed paths from r_i to c_j in G . In particular, $\text{Path}(i|j) = 0$ if $\Phi_G(i|j) = \emptyset$. Thus, the entries of Path_G belong to the \mathbb{K} -algebra \mathcal{L}_G of Laurent polynomials generated by the set W of inner vertices of G subject to (G1)–(G3).

Flows. Let $\mathcal{E}^{m,n}$ be the set of pairs $(I|J)$ with $I \subseteq [m]$, $J \subseteq [n]$ and $|I| = |J|$. Borrowing terminology from [5], for $(I|J) \in \mathcal{E}^{m,n}$, a set ϕ of pairwise disjoint directed paths from the source set $R_I := \{r_i : i \in I\}$ to the sink set $C_J := \{c_j : j \in J\}$ in G is called an $(I|J)$ -flow.

The set of $(I|J)$ -flows ϕ in G is denoted by $\Phi(I|J) = \Phi_G(I|J)$. We assume that the paths forming ϕ are ordered by increasing the source indices: if I consists of $i(1) < i(2) < \dots < i(k)$ and J consists of $j(1) < j(2) < \dots < j(k)$, then ℓ -th path P_ℓ in ϕ begins at $r_{i(\ell)}$, and therefore, P_ℓ ends at $c_{j(\ell)}$ (which follows from the planarity of G , the ordering of sources and sinks in the boundary of G and the fact that the paths in ϕ are disjoint). We write $\phi = (P_1, P_2, \dots, P_k)$ and (similar to path systems in [1]) define the weight of ϕ to be the ordered product

$$w(\phi) := w(P_1)w(P_2) \cdots w(P_k). \quad (2.4)$$

Generalizing a q -analog of Lindström's Lemma shown for Cauchon graphs in [1], one can express minors of path matrices via flows as follows.

Theorem 2.1 ([4]). *Let G be an (m, n) SE-graph. Then for the path matrix $\text{Path} = \text{Path}_G$ and for any $(I|J) \in \mathcal{E}^{m,n}$, there holds*

$$[I|J]_{\text{Path},q} = \sum_{\phi \in \Phi(I|J)} w(\phi). \quad (2.5)$$

An important fact is that the (quasi)commutation relations for the entries of Path_G are similar to those for the canonical generators x_{ij} of the quantum algebra \mathcal{R} in (1.1).

Proposition 2.2. *For an SE-graph G , the entries of its path matrix Path_G satisfy Manin's relations.*

(A proof, omitted here, can be given as an easy application of our flow-matching method.) This implies that the map $x_{ij} \mapsto \text{Path}_G(i|j)$ determines a homomorphism of \mathcal{R} to the subalgebra of \mathcal{L}_G generated by the entries of Path_G , i.e., Path_G is a q -matrix for any SE-graph G . A sharper property holds for the graph associated with the $m \times n$ Cauchon diagram without black cells. Namely, Corollary 3.2.5 in [2] relying on Cauchon's Algorithm [3] gives the following property (in our terms).

Theorem 2.3. *Let G be the extended grid $\Gamma_{m,n}$ (defined in Remark 1). Then Path_G is a generic q -matrix, i.e., $x_{ij} \mapsto \text{Path}_G(i|j)$ gives an injective map of \mathcal{R} to \mathcal{L}_G .*

Due to this property, the universal quadratic relations that we establish for minors of path matrices of SE-graphs turn out to be automatically valid for the algebra \mathcal{R} of quantum matrices, and vice versa.

3 Double flows, matchings, and exchange operations

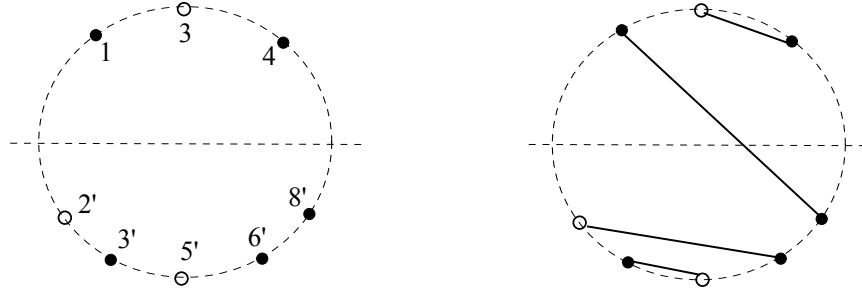
Quadratic identities of our interest involve products of the form $[I|J]_q[I'|J']_q$, where $(I|J), (I'|J') \in \mathcal{E}^{m,n}$. This leads us to a study of pairs of flows $\phi \in \Phi(I|J)$ and $\phi' \in \Phi(I'|J')$. We need some definitions and conventions, borrowing terminology from [5].

Given $I, J, I', J', \phi, \phi'$ as above, we call the pair (ϕ, ϕ') a *double flow* in G . Let

$$\begin{aligned} I^\circ &:= I - I', & J^\circ &:= J - J', & I^\bullet &:= I' - I, & J^\bullet &:= J' - J, \\ Y^r &:= I^\circ \cup I^\bullet & \text{and} & & Y^c &:= J^\circ \cup J^\bullet. \end{aligned} \quad (3.1)$$

Then $|I| = |J|$ and $|I'| = |J'|$ imply that $|I^\circ| - |I^\bullet| = |J^\circ| - |J^\bullet|$ and that $|Y^r| + |Y^c|$ is even. As before, we refer to the quadruple $(I|J, I'|J')$ as a *cortege*, and call $(I^\circ, I^\bullet, J^\circ, J^\bullet)$ the *refinement* of $(I|J, I'|J')$, or a *refined cortege*.

We interpret I° and I^\bullet as the sets of *white* and *black* elements of Y^r , respectively, and similarly for J°, J^\bullet, Y^c , and visualize these objects by use of a *circular diagram* in which the elements of Y^r (respectively Y^c) are disposed in the increasing order from left to right in the upper (respectively lower) half of a circumference O . For example, if $I^\circ = \{3\}$, $I^\bullet = \{1, 4\}$, $J^\circ = \{2', 5'\}$ and $J^\bullet = \{3', 6', 8'\}$, then the diagram is viewed as in the left fragment of the picture below. (Here, to avoid a mess, we denote the elements of Y^c with primes.)



Matchings. A partition M of $Y^r \sqcup Y^c$ into 2-element sets is called a *perfect matching* on $Y^r \sqcup Y^c$ (where \sqcup stands for the disjoint union). We say that $\pi \in M$ is: an *R-couple* if $\pi \subseteq Y^r$, a *C-couple* if $\pi \subseteq Y^c$, and an *RC-couple* if $|\pi \cap Y^r| = |\pi \cap Y^c| = 1$ (as though π “connects” two sources, two sinks, and one source and one sink, respectively). A perfect matching M is called a *feasible matching* for $(I^\circ, I^\bullet, J^\circ, J^\bullet)$ (and for $(I|J, I'|J')$) if:

(FM1) for each $\pi = \{i, j\} \in M$, the elements i, j have different colors if π is an *R-* or *C-couple*, and have the same color if π is an *RC-couple*; and

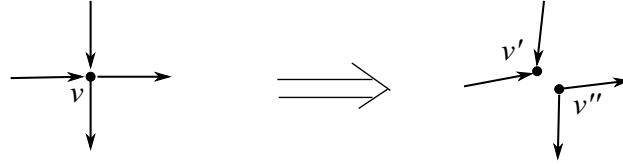
(FM2) M is *planar*, in the sense that the chords connecting the couples in the circumference O are pairwise non-intersecting.

The set of feasible matchings for $(I^\circ, I^\bullet, J^\circ, J^\bullet)$ is denoted by $\mathcal{M}_{I^\circ, I^\bullet, J^\circ, J^\bullet}$, or $\mathcal{M}(I|J, I'|J')$. One can show that this set is nonempty whenever $Y^r \sqcup Y^c \neq \emptyset$. The right fragment of the above picture illustrates an instance of feasible matchings.

Next we return to a double flow (ϕ, ϕ') as above, and our aim is to associate to it a feasible matching for $(I^\circ, I^\bullet, J^\circ, J^\bullet)$. Let V_ϕ and E_ϕ denote the sets of vertices and edges of G occurring in ϕ , respectively; and similarly for ϕ' . Consider the subgraph $\langle U \rangle$ of G induced by the set of edges

$$U := E_\phi \Delta E_{\phi'},$$

(where $A \Delta B$ denotes the symmetric difference $(A - B) \cup (B - A)$ of sets A, B). Then a vertex v of $\langle U \rangle$ has degree 1 if $v \in R_{I^\circ} \cup R_{I^\bullet} \cup C_{J^\circ} \cup C_{J^\bullet}$, and degree 2 or 4 otherwise. We modify $\langle U \rangle$ by splitting each vertex v of degree 4 in $\langle U \rangle$ into two vertices v', v'' disposed in a small neighborhood of v so that the edges entering (respectively leaving) v become entering v' (respectively leaving v''):



The resulting graph, denoted as $\langle U \rangle'$, is planar and has vertices of degree only 1 and 2. Therefore, $\langle U \rangle'$ consists of pairwise disjoint (non-directed) simple paths P'_1, \dots, P'_k and, possibly, simple cycles Q'_1, \dots, Q'_d . The corresponding images of P'_1, \dots, P'_k (respectively Q'_1, \dots, Q'_d) give paths P_1, \dots, P_k (respectively cycles Q_1, \dots, Q_d) in $\langle U \rangle$. When $\langle U \rangle$ has vertices of degree 4, some of the latter paths and cycles may be self-intersecting and may “touch”, but not “cross”, each other. It is not difficult to see the following

Lemma 3.1. (i) $k = (|I^\circ| + |I^\bullet| + |J^\circ| + |J^\bullet|)/2$;

(ii) the set of endvertices of P_1, \dots, P_k is $R_{I^\circ \cup I^\bullet} \cup C_{J^\circ \cup J^\bullet}$; moreover, each P_i connects either R_{I° and R_{I^\bullet} , or C_{J° and C_{J^\bullet} , or R_{I° and C_{J° , or R_{I^\bullet} and C_{J^\bullet} ;

(iii) in each path P_i , the edges of ϕ and the edges of ϕ' have different directions (say, the former edges are all forward, and the latter ones are all backward).

Thus, each P_i is representable as a concatenation $P_i^{(1)} \circ P_i^{(2)} \circ \dots \circ P_i^{(\ell)}$ of forwardly and backwardly directed paths which are alternately contained in ϕ and ϕ' , called the *segments* of P_i . We say that P_i is an *exchange path*. The endvertices of P_i determine a pair of elements of $Y^r \sqcup Y^c$, denoted by π_i . Then $M := \{\pi_1, \dots, \pi_k\}$ is a perfect matching on $Y^r \sqcup Y^c$. Moreover, it is feasible, since (FM1) follows from Lemma 3.1(ii), and (FM2) from the fact that P'_1, \dots, P'_k are disjoint simple paths in $\langle U \rangle'$. We denote M as $M(\phi, \phi')$, and for $\pi \in M$, denote the exchange path P_i corresponding to π by $P(\pi)$.

Corollary 3.2. $M(\phi, \phi') \in \mathcal{M}_{I^\circ, I^\bullet, J^\circ, J^\bullet}$.

Flow exchange operation. It rearranges a given double flow (ϕ, ϕ') for $(I|J, I'|J')$ into another double flow (ψ, ψ') for some $(\tilde{I}|\tilde{J}, \tilde{I}'|\tilde{J}')$, as follows. Fix a submatching $\Pi \subseteq M(\phi, \phi')$, and combine the exchange paths related to Π , forming the set of edges

$$\mathcal{E} := \cup(E_{P(\pi)} : \pi \in \Pi).$$

(where $E_{P(\pi)}$ is the edge set of $P(\pi)$). Using [Lemma 3.1](#), one can show the following

Lemma 3.3. *Let $V_\Pi := \cup(\pi \in \Pi)$. Define*

$$\tilde{I} := I \Delta (V_\Pi \cap Y^r), \quad \tilde{I}' := I' \Delta (V_\Pi \cap Y^r), \quad \tilde{J} := J \Delta (V_\Pi \cap Y^c), \quad \tilde{J}' := J' \Delta (V_\Pi \cap Y^c).$$

Then the subgraph ψ induced by $E_\phi \Delta \mathcal{E}$ gives a $(\tilde{I}|\tilde{J})$ -flow, and the subgraph ψ' induced by $E_{\phi'} \Delta \mathcal{E}$ gives a $(\tilde{I}'|\tilde{J}')$ -flow in G . Furthermore, $E_\psi \cup E_{\psi'} = E_\phi \cup E_{\phi'}$, $E_\psi \Delta E_{\psi'} = E_\phi \Delta E_{\phi'}$ ($= U$), and $M(\psi, \psi') = M(\phi, \phi')$.

We call the transformation $(\phi, \phi') \xrightarrow{\Pi} (\psi, \psi')$ in this lemma the *flow exchange operation* for (ϕ, ϕ') using $\Pi \subseteq M(\phi, \phi')$. Clearly the exchange operation applied to (ψ, ψ') using the same Π returns (ϕ, ϕ') .

So far our description has been close to that given for the commutative case in [5]. From now on we will essentially deal with the quantum version. The next theorem serves the main working tool in our arguments; its proof, based on a bulk of combinatorics on paths and flows, is given in [4].

Theorem 3.4. *Let ϕ be an $(I|J)$ -flow, and ϕ' an $(I'|J')$ -flow in G . Let (ψ, ψ') be the double flow obtained from (ϕ, ϕ') by the flow exchange operation using a single couple $\pi = \{i, j\} \in M(\phi, \phi')$. Then:*

(i) *when π is an R- or C-couple and $i < j$,*

$$\begin{aligned} w(\phi)w(\phi') &= qw(\psi)w(\psi') && \text{in case } i \in I \cup J; \\ w(\phi)w(\phi') &= q^{-1}w(\psi)w(\psi') && \text{in case } i \in I' \cup J'; \end{aligned}$$

(ii) *when π is an RC-couple, $w(\phi)w(\phi') = w(\psi)w(\psi')$.*

An immediate consequence from this theorem is the following.

Corollary 3.5. *For an $(I|J)$ -flow ϕ and an $(I'|J')$ -flow ϕ' , let (ψ, ψ') be obtained from (ϕ, ϕ') by the flow exchange operation using a set $\Pi \subseteq M(\phi, \phi')$. Then*

$$w(\phi)w(\phi') = q^{\zeta^\circ - \zeta^\bullet} w(\psi)w(\psi'), \tag{3.2}$$

where $\zeta^\circ = \zeta^\circ(I|J, I'|J'; \Pi)$ (respectively $\zeta^\bullet = \zeta^\bullet(I|J, I'|J'; \Pi)$) is the amount of R- or C-couples $\pi = \{i, j\} \in \Pi$ such that $i < j$ and $i \in I \cup J$ (respectively $i \in I' \cup J'$).

Indeed, the flow exchange operation using the whole Π reduces to performing, step by step, the exchange operations using single couples $\pi \in \Pi$ (taking into account that for any current double flow (η, η') occurring in the process, the sets $E_\eta \cup E_{\eta'}$ and $E_\eta \Delta E_{\eta'}$, as well as the matching $M(\eta, \eta')$, do not change; cf. [Lemma 3.3](#)). Then (3.2) follows from [Theorem 3.4](#).

4 Quadratic identities and the q -balancedness

As before, we consider an (m, n) SE-graph $G = (V, E; R, C)$ and deal with q -minors $[I|J] = [I|J]_{\text{Path}, q}$ of its path matrix $\text{Path} = \text{Path}_G$. In this section, based on [Corollary 3.5](#) and developing a quantum version of the flow-matching method elaborated for the commutative case in [5], we establish sufficient conditions of a general form on quadratic relations for q -minors to be valid independently of G and some other data (mentioned in Remark 2 below), referring to them as “universal quadratic identities”.

Relations that we deal with are of the form

$$\sum_{\mathcal{I}} q^{\alpha(I|J, I'|J')} [I|J][I'|J'] = \sum_{\mathcal{K}} q^{\beta(K|L, K'|L')} [K|L][K'|L'], \quad (4.1)$$

where α, β are integer-valued, \mathcal{I} is a family of corteges $(I|J, I'|J') \in \mathcal{E}^{m,n} \times \mathcal{E}^{m,n}$ (with possible multiplicities), and similarly for \mathcal{K} . (Then (4.1) is equivalent to (1.2) with δ corresponding to (α, β) .) We assume that \mathcal{I} and \mathcal{K} are *homogeneous*, in the sense that for any $(I|J, I'|J') \in \mathcal{I}$ and $(K|L, K'|L') \in \mathcal{K}$,

$$I \cup I' = K \cup K', \quad J \cup J' = L \cup L', \quad I \cap I' = K \cap K', \quad J \cap J' = L \cap L'. \quad (4.2)$$

Moreover, we shall see that only the refinements $(I^\circ, I^\bullet, J^\circ, J^\bullet)$ and $(K^\circ, K^\bullet, L^\circ, L^\bullet)$ are important, whereas the sets $I \cap I'$ and $J \cap J'$ are, in fact, indifferent.

To formulate our validity criterion, we need some definitions and notation.

- A tuple $(I|J, I'|J'; M)$, where $(I|J, I'|J') \in \mathcal{I}$ and $M \in \mathcal{M}_{I^\circ, I^\bullet, J^\circ, J^\bullet}$ (cf. (FM1)–(FM2)), is called a *configuration* for \mathcal{I} . The family of all configurations for \mathcal{I} is denoted by $\mathbf{C}(\mathcal{I})$. Similarly, we define the family $\mathbf{C}(\mathcal{K})$ of configurations for \mathcal{K} .

- Define $\mathbf{M}(\mathcal{I})$ to be the family of all matchings M (with possible multiplicities) occurring in the members of $\mathbf{C}(\mathcal{I})$. Define $\mathbf{M}(\mathcal{K})$ in a similar way.

- Families \mathcal{I} and \mathcal{K} are called *balanced* (borrowing terminology from [5]) if there exists a bijection $(I|J, I'|J'; M) \mapsto (K|K', L|L'; M')$ between $\mathbf{C}(\mathcal{I})$ and $\mathbf{C}(\mathcal{K})$ such that $M = M'$. In other words, \mathcal{I} and \mathcal{K} are balanced if $\mathbf{M}(\mathcal{I}) = \mathbf{M}(\mathcal{K})$.

- Families \mathcal{I} and \mathcal{K} along with functions $\alpha : \mathcal{I} \rightarrow \mathbb{Z}$ and $\beta : \mathcal{K} \rightarrow \mathbb{Z}$ are called *q -balanced* if there exists a bijection γ as above such that, for each $(I|J, I'|J'; M) \in \mathbf{C}(\mathcal{I})$ and for $(K|K', L|L'; M) = \gamma(I|J, I'|J'; M)$, there holds

$$\beta(K|K', L|L') - \alpha(I|J, I'|J') = \zeta^\circ - \zeta^\bullet. \quad (4.3)$$

(In particular, \mathcal{I}, \mathcal{K} are balanced.) Here $\zeta^\circ, \zeta^\bullet$ are defined according to [Corollary 3.5](#). Namely, $\zeta^\circ = \zeta^\circ(I|J, I'|J'; \Pi)$ and $\zeta^\bullet = \zeta^\bullet(I|J, I'|J'; \Pi)$, where Π is the set of couples $\pi \in M$ such that the colorings of π in the refined corteges $(I^\circ, I^\bullet, J^\circ, J^\bullet)$ and $(K^\circ, K^\bullet, L^\circ, L^\bullet)$ are different. (Then ζ° (ζ^\bullet) is the number of R - and C -couples $\{i, j\} \in \Pi$ with $i < j$ and $i \in I^\circ \cup J^\circ$ (respectively $i \in I^\bullet \cup J^\bullet$.) We say that $(K^\circ, K^\bullet, L^\circ, L^\bullet)$ is obtained from $(I^\circ, I^\bullet, J^\circ, J^\bullet)$ by the *index exchange operation* using Π .

Theorem 4.1. *Let \mathcal{I} and \mathcal{K} be homogeneous families on $\mathcal{E}^{m,n} \times \mathcal{E}^{m,n}$, and let $\alpha : \mathcal{I} \rightarrow \mathbb{Z}$ and $\beta : \mathcal{K} \rightarrow \mathbb{Z}$. Suppose that $\mathcal{I}, \mathcal{K}, \alpha, \beta$ are q -balanced. Then (4.1) is valid for q -minors of the path matrix of any (m, n) SE-graph $G = (V, E; R, C)$.*

Proof. We fix G and denote by $\mathcal{D}(I|J, I'|J')$ the set of double flows for $(I|J, I'|J') \in \mathcal{I} \cup \mathcal{K}$ in G . A summand concerning $(I|J, I'|J') \in \mathcal{I}$ in the left side of (4.1) can be expressed via double flows as follows, ignoring the factor of $q^{\alpha(\cdot)}$:

$$\begin{aligned} [I|J][I'|J'] &= \left(\sum_{\phi \in \Phi_G(I|J)} w(\phi) \right) \times \left(\sum_{\phi' \in \Phi_G(I'|J')} w(\phi') \right) \\ &= \sum_{(\phi, \phi') \in \mathcal{D}(I|J, I'|J')} w(\phi)w(\phi') \\ &= \sum_{M \in \mathcal{M}_{I^\circ, I', J^\circ, J'}} \sum_{(\phi, \phi') \in \mathcal{D}(I|J, I'|J') : M(\phi, \phi') = M} w(\phi)w(\phi'). \end{aligned} \quad (4.4)$$

The summand for $(K|L, K'|L') \in \mathcal{K}$ in the right side of (4.1) is expressed similarly.

Consider a configuration $S = (I|J, I'|J'; M) \in \mathbf{C}(\mathcal{I})$ and suppose that (ϕ, ϕ') is a double flow for $(I|J, I'|J')$ with $M(\phi, \phi') = M$ (if such a double flow in G exists). Since $\mathcal{I}, \mathcal{K}, \alpha, \beta$ are q -balanced, S is bijective to some configuration $S' = (K|L, K'|L'; M) \in \mathbf{C}(\mathcal{K})$ satisfying (4.3). As explained earlier, the cortege $(K|L, K'|L')$ is obtained from $(I|J, I'|J')$ by the index exchange operation using some $\Pi \subseteq M$. Then the flow exchange operation applied to (ϕ, ϕ') using this Π results in a double flow (ψ, ψ') for $(K|L, K'|L')$ which satisfies relation (3.2) in Corollary 3.5. Comparing (3.2) with (4.3), we observe that

$$q^{\alpha(I|J, I'|J')} w(\phi)w(\phi') = q^{\beta(K|L, K'|L')} w(\psi)w(\psi').$$

Furthermore, such a map $(\phi, \phi') \mapsto (\psi, \psi')$ gives a bijection between all double flows concerning configurations in $\mathbf{C}(\mathcal{I})$ and those in $\mathbf{C}(\mathcal{K})$. Now the desired equality (4.1) follows by considering the last term in expression (4.4) and the corresponding term in the analogous expression concerning \mathcal{K} . \square

As a consequence of Theorems 2.3 and 4.1, the following result is obtained.

Corollary 4.2. *If $\mathcal{I}, \mathcal{K}, \alpha, \beta$ as above are q -balanced, then relation (4.1) is valid for the corresponding minors in the algebra \mathcal{R} of quantum $m \times n$ matrices.*

Remark 2. When speaking of a *universal quadratic identity* of the form (4.1) with homogeneous \mathcal{I} and \mathcal{K} , abbreviated as a *UQ identity*, we mean that it depends neither on the SE-graph G nor on the field \mathbb{K} and element $q \in \mathbb{K}^*$, and that the index sets can be modified as follows. Given $(I|J, I'|J') \in \mathcal{I}$, let $A := I \Delta I'$, $B := J \Delta J'$, $S := I \cap I'$ and $T := J \cap J'$. Take arbitrary $\tilde{m} \geq |A|$ and $\tilde{n} \geq |B|$, disjoint sets $\tilde{A}, \tilde{S} \subseteq [\tilde{m}]$, and disjoint sets $\tilde{B}, \tilde{T} \subseteq [\tilde{n}]$ with $|\tilde{A}| = |A|$, $|\tilde{B}| = |B|$, $|\tilde{S}| - |\tilde{T}| = |S| - |T|$. Let $\lambda : A \rightarrow \tilde{A}$ and $\mu : B \rightarrow \tilde{B}$ be the order preserving maps. Transform each $(I|J, I'|J') \in \mathcal{I}$ into $(\tilde{I}|\tilde{J}, \tilde{I}'|\tilde{J}')$, obtaining a new family $\tilde{\mathcal{I}}$ on $\mathcal{E}^{\tilde{m}, \tilde{n}} \times \mathcal{E}^{\tilde{m}, \tilde{n}}$, where

$$\tilde{I} := \tilde{S} \cup \lambda(I - S), \quad \tilde{I}' := \tilde{S} \cup \lambda(I' - S), \quad \tilde{J} := \tilde{T} \cup \mu(J - T), \quad \tilde{J}' := \tilde{T} \cup \mu(J' - T).$$

Turn \mathcal{K} into $\tilde{\mathcal{K}}$ in a similar way. One can see that if $\mathcal{I}, \mathcal{K}, \alpha, \beta$ are q -balanced, then so are $\tilde{\mathcal{I}}, \tilde{\mathcal{K}}$, keeping α, β . Therefore, if (4.1) is valid for \mathcal{I}, \mathcal{K} , then it is valid for $\tilde{\mathcal{I}}, \tilde{\mathcal{K}}$ as well.

One can say that identity (4.1), where all summands have positive signs, is written in the canonical form. Sometimes, however, it is more convenient to consider equivalent identities having negative summands in one or both sides (e.g., of the form (1.2)).

We can suggest a rather simple algorithm which has as the input a corresponding quadruple $\mathcal{I}, \mathcal{K}, \alpha, \beta$ and recognizes the q -balanced for it. Therefore, in light of **Theorems 4.1** and **4.3**, the algorithm decides whether or not the given quadruple determines a UQ identity of the form (4.1).

Algorithm. Compute the set $\mathcal{M}_{I^\circ, I^\bullet, J^\circ, J^\bullet}$ of feasible matchings M for each $(I|J, I'|J') \in \mathcal{I}$, and similarly for \mathcal{K} . For each instance M occurring there, extract the family $\mathbf{C}_M(\mathcal{I})$ of all configurations concerning M in $\mathbf{C}(\mathcal{I})$, and extract a similar family $\mathbf{C}_M(\mathcal{K})$ in $\mathbf{C}(\mathcal{K})$. If $|\mathbf{C}_M(\mathcal{I})| \neq |\mathbf{C}_M(\mathcal{K})|$ for at least one instance M , then \mathcal{I} and \mathcal{K} are not balanced at all. Otherwise for each M , we seek for a required bijection $\gamma_M : \mathbf{C}_M(\mathcal{I}) \rightarrow \mathbf{C}_M(\mathcal{K})$ by solving the maximum matching problem in the corresponding bipartite graph H_M . More precisely, the vertices of H_M are the tuples $(I|J, I'|J'; M)$ and $(K|L, K'|L'; M)$ occurring in $\mathbf{C}_M(\mathcal{I})$ and $\mathbf{C}_M(\mathcal{K})$, and such tuples are connected by edge in H_M if they obey (4.3). Find a maximum matching N in H_M . If $|N| = |\mathbf{C}_M(\mathcal{I})|$, then N determines the desired γ_M in a natural way. Taking together, these γ_M give a bijection between $\mathbf{C}(\mathcal{I})$ and $\mathbf{C}(\mathcal{K})$ as required, implying that $\mathcal{I}, \mathcal{K}, \alpha, \beta$ are q -balanced. And if $|N| < |\mathbf{C}_M(\mathcal{I})|$ for at least one instance M , then the algorithm declares the non- q -balancedness.

Example. Next we illustrate our method with one example. Recall that sets $I, J \subseteq [n]$ are called *weakly separated* if, up to renaming I and J , there holds (*): $|I| \geq |J|$, and $J - I$ has a partition $\{J_1, J_2\}$ such that $J_1 < I - J < J_2$ (where we write $X < Y$ if $x < y$ for any $x \in X$ and $y \in Y$). Let $k := |I|$ and $\ell := |J|$. Leclerc and Zelevinsky [9] proved that: *two quantum flag minors* $[I] := [[k|I]_q$ and $[J] := [[\ell|J]_q$ quasicommute, i.e., satisfy $[I][J] = q^c [J][I]$ for some $c \in \mathbb{Z}$, if and only if I and J are weakly separated. Moreover, assuming (*), c is equal to $|J_2| - |J_1|$. One can show “if” part of this theorem as follows. (For “only if” part, a characterization of quasicommuting non-flag minors, and many other applications of our method, see arXiv:1604.00338[math.QA].)

Let $A := [k] - [\ell]$. Assuming (*), one can see that $\mathcal{M}([k|I, [\ell|J])$ has exactly one feasible matching M ; namely, J_1 is coupled with the first $|J_1|$ elements of $I - J$, J_2 is coupled with the last $|J_2|$ elements of $I - J$ (forming all C-couples), and the rest of $I - J$ is coupled with A (forming all RC-couples). Observe that the index exchange operation applied to the cortege $([k|I, [\ell|J])$ using the whole M swaps $([k|I)$ and $([\ell|J)$ (as it changes the colors of all elements of $(I - J) \cup (J - I) \cup A$). Also the set of C-couples of M consist of $|J_1|$ couples $\{i, j\}$ with $i < j$ and $i \in J_1$, and $|J_2|$ couples $\{i, j\}$ with $i < j$ and $j \in J_2$. This gives $\zeta^\circ = |J_2|$ and $\zeta^\bullet = |J_1|$. Hence the (one-element) families $\mathcal{I} = \{([k|I, [\ell|J])\}$ and $\mathcal{K} = \{([\ell|J, [k|I])\}$ along with $\alpha(([k|I, [\ell|J])) = 0$ and

$\beta(([\ell|J, [k|I])) = |J_2| - |J_1|$ are q -balanced. Now “if” part (with $c = |J_2| - |J_1|$) follows from [Theorem 4.1](#).

Finally, we formulate (without a proof) a converse assertion to [Theorem 4.1](#), saying that the q -balancedness condition is necessary as well. This gives a complete characterization for the UQ identities on quantum minors.

Theorem 4.3. *Let \mathbb{K} be a field of characteristic zero and let $q \in \mathbb{K}^*$ be transcendental over \mathbb{Q} . Suppose that $\mathcal{I}, \mathcal{K}, \alpha, \beta$ (as in [Section 4](#)) are not q -balanced. Then there exists, and can be explicitly constructed, an SE-graph G for which relation [\(4.1\)](#) is violated.*

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